

# **Cyclotomic Rings, Lattices and Space-Time Block Codes**

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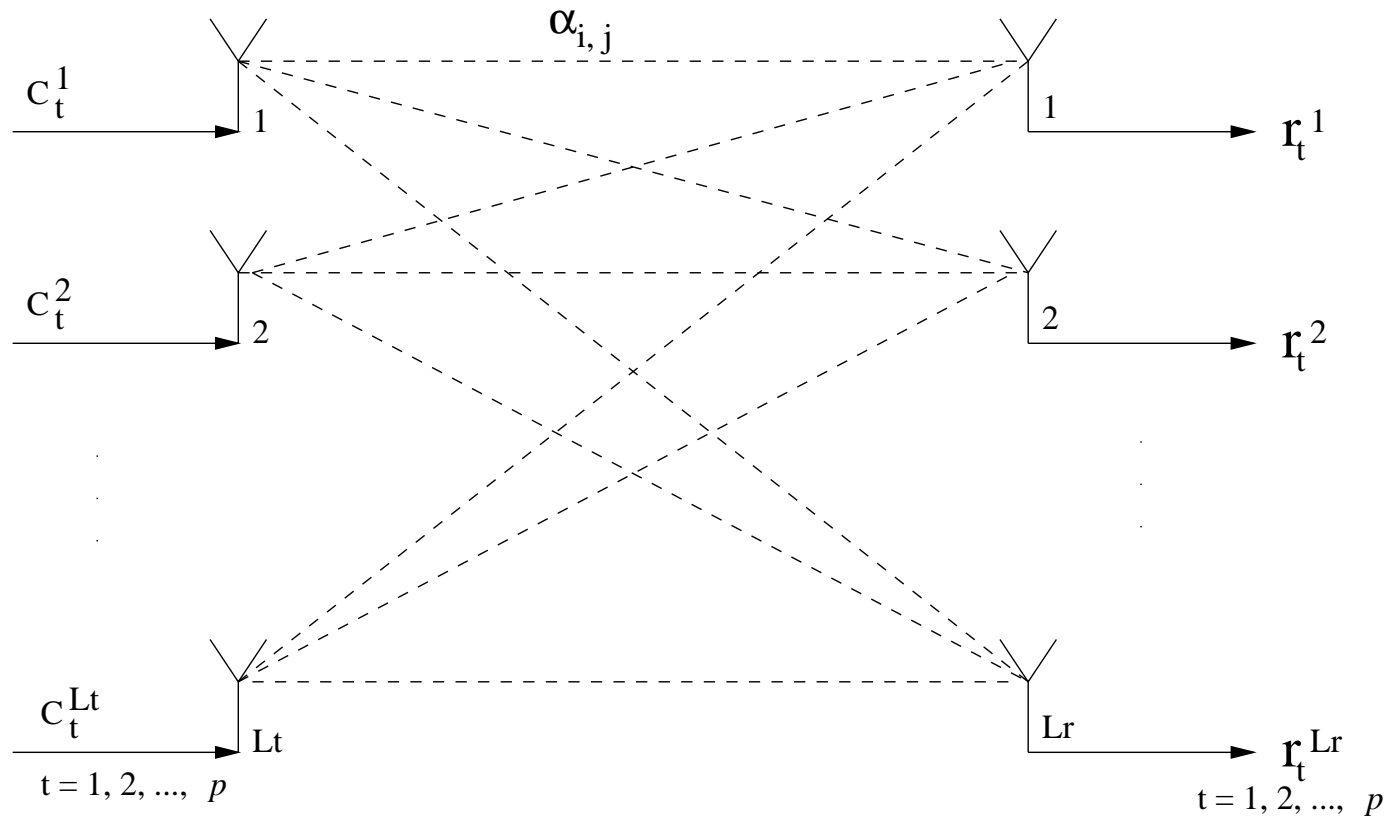
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# OUTLINE

- Motivation
- Cyclotomic Rings
- Cyclotomic Lattices
  - Complex Lattices and Some Existing Examples
  - Cyclotomic Lattices of Full Diversity
  - Cyclotomic Diagonal Space-Time Codes
- Optimal Cyclotomic Lattices and Diagonal Space-Time Codes
- Some Simulation Results
- Conclusion and Future Research

# Multiple Antenna System



$L_t$  transmit antennas;  $L_r$  receive antennas;  $p$  is the time block size.

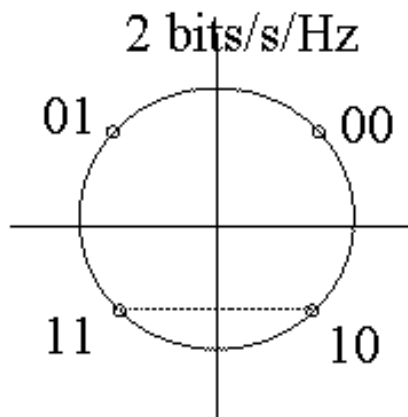
$\alpha_{ij}$  is the channel coefficient from  $i$ th transmit to  $j$ th receive antenna and is a random variable.

## Advantage of Multi-Antenna System: Capacity Gain

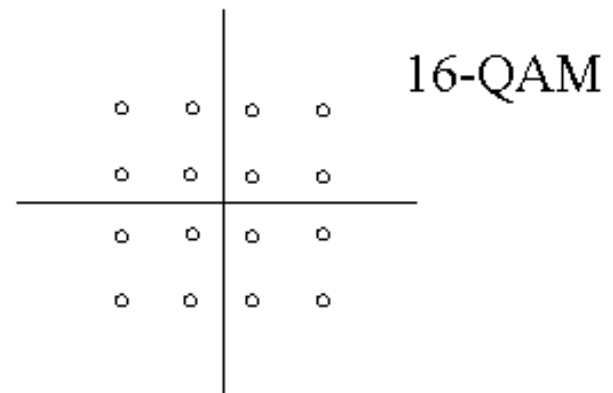
- Teletar (1995), Foschini and Gans (1998) proved that the capacity of a multi-antenna system is proportional to  $\min\{L_t, L_r\}$ .
  - Theoretically, the more transmit and receive antennas, the better the capacity!
  - Practically, how can we achieve the capacity (gain)?
  - Shannon communication theory tells us that the capacity can be achieved by coding and modulation, **BUT HOW??**
    - **One of the most active research areas in communications theory!**

## Single Antenna Coding and Modulation

- Low rate transmission: 1 bit modulated to 1 number/symbol (BPSK)
- High rate transmission: multiple bits modulated to 1 number/symbol (M-QAM, M-PSK)
- Consider 4, 16-QAM: 2, 4 bits become a complex number/symbol:



These 4 points are optimal:  
The minimum distance is maximal



These 16 points are almost optimal

## What is Multiple Antenna Coding and Modulation: Signal Model

- Transmit and receive signal model:

$$Y = AC + W,$$

where

$Y = (r_t^i)_{L_r \times p}$ : Received signal matrix

$A = (\alpha_{ij})_{L_r \times L_t}$ : Channel coefficient matrix

$C = (c_t^j)_{L_t \times p}$ : Transmit signal matrix

$W = (w_{i,t})_{L_r \times p}$ : Additive white Gaussian noise (AWGN) matrix.

## What is Multiple Antenna Coding and Modulation

- Multiple antenna coding/modulation: Binary information bits are modulated/mapped into  $L_t \times p$  matrices and these matrices are taken from a pre-designed matrix set called *space-time code*.
- How to design a space-time code: It should be designed in such a way that the error probability at the receiver is minimized.

## Space-Time Code Design Criteria

Based on the pair-wise error probability from the maximum-likelihood (ML) decoding, Guey-Fitz-Bell-Kuo and Tarokh-Seshadri-Calderbank proposed the following rank and diversity product criteria:

- **Rank criterion:** Any difference matrix of any two distinct matrices in a space-time code  $\mathcal{C}$  has full rank;
- **Diversity product criterion** (or coding advantage or product distance):

$$\xi(\mathcal{C}) = \min_{C \neq \tilde{C} \in \mathcal{C}} |\det((C - \tilde{C})^\dagger (C - \tilde{C}))|$$

is as large as possible.



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## Some Existing Space-Time Coding Schemes

- Space-time block codes
  - BLAST–No full diversity
  - Alamouti scheme-Orthogonal space-time block codes from orthogonal designs, quasi orthogonal space-time codes– Fast decoding with good performance BUT symbol rates are limited!
  - Unitary space-time codes: Group codes, Caley transforms, parametric codes, packing theory etc. – Good performance but no fast decoding in general, no systematic designs!
  - Space-time codes from binary linear codes – Performance is limited!
  - **Linear lattice based codes** – Fast sphere decoding, high rates, systematic: **My focus**
- Space-time trellis codes

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# Linear Lattice Based Space-Time Codes and Motivation

- BLAST schemes
- (Quasi) Orthogonal space-time codes (Alamouti, Tarokh et al, Jafarkhani, Tirkkonen et al etc.)
- Linear dispersive codes (Hassibi-Hochwald, Heath et al, Sandhu et al, etc.)
- Signal Space Diversity Codes and Diagonal Space-Time Codes using Algebraic Number Theory (Boutros-Viterbo, Giraud-Boutillon-Belfiore, Damen-Meraim-Belfiore, Sethuraman-Rajan, .....
  - The existing lattice based space-time block codes
    - \* Not **concrete** but simply some abstract algebraic numbers
    - \* Mostly based on square QAM, i.e., square lattice.
    - \* Not optimized in terms of diversity product

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- The **focus** of this presentation:
    - \* Propose a general systematic and **concrete** cyclotomic full diversity lattices.
    - \* Propose optimal cyclotomic lattices and space-time block codes in terms of maximized diversity product for fixed mean signal power.

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## General Problem Description

- Let  $L_t$  be the number of transmit antennas.
- Let  $G$  be an  $L_t \times L_t$  matrix and

$$[\mathbf{y}_1, \dots, \mathbf{y}_{L_t}]^T = G[\mathbf{x}_1, \dots, \mathbf{x}_{L_t}]^T,$$

where  $\mathbf{x}_i$  are information symbols.

- A diagonal space-time code  $\Omega$  consists of  $L_t \times L_t$  matrices of the form  $\text{diag}(\mathbf{y}_1, \dots, \mathbf{y}_{L_t})$ .

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- We are interested in such a diagonal space-time code  $\Omega$  that
    - (i) it has the full rank property, i.e., any difference matrix of any two distinct matrices in  $\Omega$  has full rank; and
    - (ii) its following diversity product is as large as possible:

$$\xi = \min_{\text{diag}(\mathbf{y}_1, \dots, \mathbf{y}_{L_t}) \neq \text{diag}(\mathbf{e}_1, \dots, \mathbf{e}_{L_t}) \in \Omega} \prod_{i=1}^{L_t} |\mathbf{y}_i - \mathbf{e}_i|^2,$$

where the transmission signal mean power of  $\mathbf{y}_i$  is fixed, or **equivalently**, the transmission signal mean power is minimized, when the diversity product is fixed.

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## Cyclotomic Rings

- Let  $\zeta_m = \exp(\mathbf{j} \frac{2\pi}{m})$  be the  $m$ th root of unity.
- Let  $\mathbb{Z}[\zeta_m]$  denote the ring generated by  $\mathbb{Z}$ , all integers, and  $\zeta_m$ . It is called a *cyclotomic ring*.
- When  $m = 4$ ,  $\mathbb{Z}[\zeta_m] = \mathbb{Z}[\mathbf{j}]$  that is called Gaussian integers.
- When  $m = 3, 6$ ,  $\mathbb{Z}[\zeta_3] = \mathbb{Z}[\zeta_6]$  that is called Eisenstein integers.
- An important result from algebraic number theory: For a fixed  $L$ ,

$$\min_{(0, \dots, 0) \neq (\mathbf{x}_1, \dots, \mathbf{x}_L) \in (\mathbb{Z}[\zeta_m])^L} \prod_{i=1}^L |\mathbf{x}_i| = 1$$

where  $m = 3, 4, 6$ , i.e., for Gaussian or Eisenstein integers.

- We next want to define **cyclotomic lattices**.

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## Real Lattices

- An  $n$ -dimensional *real lattice*  $\Lambda_n(K)$  is a subset in  $\mathbb{R}^n$ :

$$\Lambda_n(K) = \left\{ \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix} = K \begin{bmatrix} z_1 \\ \vdots \\ z_n \end{bmatrix} \mid z_i \in \mathbb{Z} \text{ for } 1 \leq i \leq n \right\},$$

where  $\mathbb{Z}$  is the ring of all integers and  $K$  is an  $n \times n$  real matrix of full rank and called the generating matrix of the real lattice  $\Lambda_n(K)$  and  $\det(\Lambda_n(K)) \triangleq |\det(K)|$ .

- Every point  $[x_1, x_2]^T$  in a two dimensional real lattice  $\Lambda_2(K)$  is treated equivalently as a complex number  $\mathbf{x} = x_1 + jx_2$  in the complex plane  $\mathbb{C}$ .
- For  $\zeta_m = \exp(j \frac{2\pi}{m})$ , we use  $\Lambda_{\zeta_m}$  to denote the two dimensional real lattice

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with the generating matrix

$$K_{\zeta_m} = \begin{bmatrix} 1 & \cos(\frac{2\pi}{m}) \\ 0 & \sin(\frac{2\pi}{m}) \end{bmatrix} = \begin{bmatrix} 1 & \operatorname{Re}(\zeta_m) \\ 0 & \operatorname{Im}(\zeta_m) \end{bmatrix}.$$

Thus,  $\Lambda_{\zeta_m} = \Lambda_2(K_{\zeta_m})$ .

- $\Lambda_{\zeta_m} \subset \mathbb{Z}[\zeta_m]$ ,  $\Lambda_{\zeta_4} = \mathbb{Z}[\zeta_4] = \mathbb{Z}[j]$ , and  $\Lambda_{\zeta_3} = \Lambda_{\zeta_6} = \mathbb{Z}[\zeta_3] = \mathbb{Z}[\zeta_6]$



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## Complex Lattices

- **Definition:** An  $n$ -dimensional complex lattice  $\Gamma_n(G)$  over a two dimensional real lattice  $\Lambda_2(K)$  is a subset of  $\mathbb{C}^n$ :

$$\Gamma_n(G) = \left\{ \begin{bmatrix} \mathbf{y}_1 \\ \vdots \\ \mathbf{y}_n \end{bmatrix} = G \begin{bmatrix} \mathbf{x}_1 \\ \vdots \\ \mathbf{x}_n \end{bmatrix} \mid \mathbf{x}_i \in \Lambda_2(K), \text{ for } 1 \leq i \leq n \right\},$$

where  $G$  is an  $n \times n$  complex matrix of full rank and called the generating matrix of the complex lattice  $\Gamma_n(G)$ . The above complex lattice is called a full diversity lattice if it satisfies

$$\sum_{i=1}^n |\mathbf{y}_i|^2 > 0$$

for any non-zero vector  $[\mathbf{x}_1, \dots, \mathbf{x}_n]^T \neq [0, \dots, 0]^T$  in  $(\Lambda_2(K))^n$ .

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- **Examples**

- Rotated Codes Based on QAM on the Square Lattice:  $G_2$  and  $G_4$ , 2-point and 4-point DFT matrices based on  $\mathbb{Z}[\zeta_4]$  (Boutros-Viterbo (98) and Giraud-Boutillon-Belfiore (97))
- Good Codes for Fading Channels as well as Gaussian Channels:  $D_4$ ,  $E_6$ ,  $E_8$ ,  $G_2$ ,  $G_3$ ,  $G_4$  based on  $\mathbb{Z}[\zeta_4]$  and  $\mathbb{Z}[\zeta_3]$  (Boutros-Viterbo (98) and Giraud-Boutillon-Belfiore (97))
- Diagonal Algebraic Space-Time Block Codes – DAST Block Codes based on  $\mathbb{Z}[\zeta_4]$ :  $M_n$  (Boutros-Viterbo (98) and Damen-Meraim-Belfiore (02))

# Cyclotomic Lattices

- For two positive integers  $n$  and  $m$ , let  $N = mn$  and

$$L_t = \frac{\phi(N)}{\phi(m)},$$

where  $\phi(N)$  and  $\phi(m)$  are the Euler numbers of  $N$  and  $m$ , respectively.

- Then, there are total  $L_t$  distinct integers  $n_i$ ,  $1 \leq i \leq L_t$ , with  $0 = n_1 < n_2 < \dots < n_{L_t} \leq n - 1$  such that  $1 + n_i m$  and  $N$  are co-prime for any  $1 \leq i \leq L_t$ .
- With these  $L_t$  integers, we define

$$G_{m,n} \triangleq \begin{bmatrix} \zeta_N & \zeta_N^2 & \dots & \zeta_N^{L_t} \\ \zeta_N^{1+n_2 m} & \zeta_N^{2(1+n_2 m)} & \dots & \zeta_N^{L_t(1+n_2 m)} \\ \vdots & \vdots & \ddots & \vdots \\ \zeta_N^{1+n_{L_t} m} & \zeta_N^{2(1+n_{L_t} m)} & \dots & \zeta_N^{L_t(1+n_{L_t} m)} \end{bmatrix}_{L_t \times L_t},$$

where  $\zeta_N = \exp(j \frac{2\pi}{N})$ .

- **Definition:** An  $L_t$  dimensional complex lattice  $\Gamma_{L_t}(G_{m,n})$  over  $\Lambda_{\zeta_m}$  is called a cyclotomic lattice, where  $G_{m,n}$  is defined above and  $\Lambda_{\zeta_m}$  is the two dimensional real lattice with the generating matrix  $K_{\zeta_m}$ . Its minimum product  $d_{\min}(\Gamma_{L_t}(G_{m,n}))$  is defined by

$$d_{\min}(\Gamma_{L_t}(G_{m,n})) \triangleq \min_{[0, \dots, 0]^T \neq [\mathbf{y}_1, \dots, \mathbf{y}_{L_t}]^T \in \Gamma_{L_t}(G_{m,n})} \left| \sum_{i=1}^{L_t} \mathbf{y}_i \right|.$$

- Some equivalent forms of  $G_{m,n}$ :

$$G_{m,n} \triangleq \text{diag}(\zeta_N, \zeta_N^{1+n_2 m}, \dots, \zeta_N^{1+n_{L_t} m}) \hat{G}_{m,n},$$

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where

$$\hat{G}_{m,n} \triangleq \begin{bmatrix} 1 & \zeta_N & \cdots & \zeta_N^{L_t-1} \\ 1 & \zeta_N^{1+n_2m} & \cdots & \zeta_N^{(L_t-1)(1+n_2m)} \\ \vdots & \vdots & \ddots & \vdots \\ 1 & \zeta_N^{1+n_{L_t}m} & \cdots & \zeta_N^{(L_t-1)(1+n_{L_t}m)} \end{bmatrix}_{L_t \times L_t}.$$

$$G_{m,n} = \begin{bmatrix} \zeta_n^{n_1} & \zeta_n^{2n_1} & \cdots & \zeta_n^{L_t n_1} \\ \zeta_n^{n_2} & \zeta_n^{2n_2} & \cdots & \zeta_n^{L_t n_2} \\ \vdots & \vdots & \ddots & \vdots \\ \zeta_n^{n_{L_t}} & \zeta_n^{2n_{L_t}} & \cdots & \zeta_n^{L_t n_{L_t}} \end{bmatrix}_{L_t \times L_t} \text{diag}(\zeta_N, \zeta_N^2, \cdots, \zeta_N^{L_t}).$$

- A difference with the existing results on this topic is that the above proposed cyclotomic lattice generating matrix is **concrete** and systematic.
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- **Theorem 1:** A cyclotomic lattice is a full diversity lattice.

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## Cyclotomic Diagonal Space-Time Codes

- **Definition:** A diagonal cyclotomic space-time code  $\Omega$  for  $L_t$  transmit antennas is defined by  $\Omega = \{\text{diag}(\mathbf{y}_1, \dots, \mathbf{y}_{L_t})\}$  where  $\mathbf{y}_i$  for  $1 \leq i \leq L_t$  are defined as follows:

$$[\mathbf{y}_1, \dots, \mathbf{y}_{L_t}]^T = G_{m,n}[\mathbf{x}_1, \dots, \mathbf{x}_{L_t}]$$

where  $[\mathbf{x}_1, \dots, \mathbf{x}_{L_t}]^T \in \mathcal{S} \subset (\mathbb{Z}[\zeta_m])^{L_t}$  and  $\mathcal{S}$  is a signal constellation for information symbols.

- **Theorem 2:** A diagonal cyclotomic space-time code has full diversity.
- $G_2 = \hat{G}_{4,2}$  and  $G_4 = \hat{G}_{4,4}$ .
- **Question:** For a transmit antenna number  $L_t$ , there are infinitely many cyclotomic lattices  $G_{m,n}$  from infinitely many pairs  $(m, n)$ . For a fixed  $L_t$ , which cyclotomic lattice is **optimal** in the sense that, its mean transmission power is minimized when its diversity product is fixed?

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## Optimal Cyclotomic Lattices

- The mean transmission signal power of signal points on a lattice is reciprocal to the packing density.
- From the packing theory, the packing density is reciprocal to the absolute value of the determinant of the generating matrix.
- The absolute value of the determinant of the generating matrix of a cyclotomic is

$$|\det(G_{m,n})|^2 \cdot |\det(K_{\zeta_m})|^{L_t}.$$



- 
- **Criterion:** Let  $\Gamma_{L_t}(G_{m_1, n_1})$  and  $\Gamma_{L_t}(G_{m_2, n_2})$  be two  $L_t$  dimensional cyclotomic lattices over  $\Lambda_{\zeta_{m_1}}$  and  $\Lambda_{\zeta_{m_2}}$ , respectively. We say cyclotomic lattice  $\Gamma_{L_t}(G_{m_1, n_1})$  is *better than* cyclotomic lattice  $\Gamma_{L_t}(G_{m_2, n_2})$ , written as  $\Gamma_{L_t}(G_{m_1, n_1}) \leq \Gamma_{L_t}(G_{m_2, n_2})$ , if

$$|\det(G_{m_1, n_1})| \cdot |\det(\Lambda_{\zeta_{m_1}})|^{L_t/2} \leq |\det(G_{m_2, n_2})| \cdot |\det(\Lambda_{\zeta_{m_2}})|^{L_t/2},$$

when their minimum products are the same, i.e.,

$$d_{\min}(\Gamma_{L_t}(G_{m_1, n_1})) = d_{\min}(\Gamma_{L_t}(G_{m_2, n_2})).$$

- Define the following normalized minimum product

$$\gamma_{m, n} = \frac{d_{\min}(\Lambda_{L_t}(G_{m, n}))}{|\det(\Gamma_{\zeta_m})|^{L_t/2} |\det(G_{m, n})|}$$

- 
- **Theorem 3:** If the number of transmit antennas has the form

$$L_t = \frac{\phi(3n)}{\phi(3)} \text{ or } \frac{\phi(6n)}{\phi(6)}, \text{ for some integer } n,$$

then the optimal cyclotomic lattice can be achieved by an Eisenstein cyclotomic lattice, i.e.,  $m = 3$  or  $m = 6$ , and the minimum product (or diversity product) of the optimal cyclotomic lattice is 1.

- Examples of such  $L_t$  are

$$L_t = 3^{r_1} p_2^{r_2-1} (p_2 - 1) \cdots p_k^{r_k-1} (p_k - 1),$$

where  $k \geq 1$ ,  $p_2, \dots, p_k$  are distinct primes and different from 3, and  $r_1 \geq 0, r_2 \geq 1, \dots, r_k \geq 1$  are integers, which covers 2,3,4,6,8,9,10,12,16,18,20,22,24,27,28,30,32,....

- **Corollary:** For the listed numbers,  $L_t$ , of transmit antennas, the parameters  $(m, n)$  of the optimal cyclotomic lattices  $\Gamma_{L_t}(G_{m,n})$  over  $\Lambda_{\zeta_m}$  with generating matrix  $G_{m,n}$  are listed in the following table.

$L_t$	$(m, n)$ in $G_{m,n}$	$\gamma_{m,n}$
2	$(3, 4), (4, 3), (6, 2)$	$\frac{1}{\sqrt{3}}$
3	$(3, 3), (3, 6), (6, 3)$	$\frac{1}{4.1878}$
4	$(3, 5), (3, 10), (6, 5)$	$\frac{1}{8.3852}$
6	$(3, 7), (3, 14), (6, 7)$	$\frac{1}{84.2037}$
8	$(3, 20), (4, 15), (6, 10)$	$\frac{1}{1.125 \times 10^3}$
9	$(3, 9), (3, 18), (6, 9)$	$\frac{1}{1.0303 \times 10^4}$
10	$(3, 11), (3, 22), (6, 11)$	$\frac{1}{2.3655 \times 10^4}$
12	$(3, 15), (3, 30), (6, 15)$	$\frac{1}{4.2981 \times 10^5}$
16	$(3, 40), (4, 30), (6, 20)$	$\frac{1}{3.24 \times 10^8}$
18	$(3, 21), (3, 42), (6, 21)$	$\frac{1}{1.1752 \times 10^{10}}$
20	$(3, 25), (3, 50), (6, 25)$	$\frac{1}{4.0484 \times 10^{11}}$
22	$(3, 23), (3, 46), (6, 23)$	$\frac{1}{4.083 \times 10^{13}}$
24	$(3, 35), (3, 70), (6, 35)$	$\frac{1}{9.8192 \times 10^{13}}$
27	$(3, 27), (3, 54), (6, 27)$	$\frac{1}{3.0205 \times 10^{18}}$
28	$(3, 29), (3, 58), (6, 29)$	$\frac{1}{7.3757 \times 10^{18}}$

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## **Optimal Diagonal Cyclotomic Space-Time Code Designs**

- Optimal diagonal cyclotomic space-time block codes can be designed by
  - selecting optimal cyclotomic lattices
  - selecting signal points on the optimal cyclotomic lattices with the minimum mean transmission power (i) the information symbols are independently selected (ii) the information symbols are jointly selected.
- Design Examples

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Diversity Products of Diagonal Codes for Two and Four Transmit Antennas

Bit rate b/s/Hz	Space-Time Codes, $L_t = 2$				
	$\mathbf{M}_2\text{-}\mathbb{Z}[j]\text{-QAM}$	$G_2\text{-}\mathbb{Z}[j]\text{-QAM}$	$G_2\text{-}\mathbb{Z}[j]\text{-Joint}$	$G_{6,2}\text{-}\Lambda_{\zeta_6}\text{-QAM}$	$G_{6,2}\text{-}\Lambda_{\zeta_6}\text{-Joint}$
2	$\frac{1}{4.47}$	$\frac{1}{2}$	$\frac{1}{2}$	$\frac{1}{2}$	$\frac{1}{2}$
3	$\frac{1}{5.5231}$	$\frac{1}{5}$	$\frac{1}{4.6562}$	$\frac{1}{4.3125}$	$\frac{1}{4.125}$
4	$\frac{1}{11.2}$	$\frac{1}{10}$	$\frac{1}{9.5703}$	$\frac{1}{8.75}$	$\frac{1}{8.2266}$

Bit rate b/s/Hz	Space-Time Codes, $L_t = 4$				
	$\mathbf{M}_4\text{-}\mathbb{Z}[j]\text{-QAM}$	$G_4\text{-}\mathbb{Z}[j]\text{-QAM}$	$G_4\text{-}\mathbb{Z}[j]\text{-Joint}$	$G_{6,5}\text{-}\Lambda_{\zeta_6}\text{-QAM}$	$G_{6,5}\text{-}\Lambda_{\zeta_6}\text{-Joint}$
2	$\frac{1}{640}$	$\frac{1}{256}$	$\frac{1}{94.15}$	$\frac{1}{64}$	$\frac{1}{43.0664}$
3	$\frac{1}{1000}$	$\frac{1}{400}$	$\frac{1}{323.2265}$	$\frac{1}{297.5625}$	$\frac{1}{170.514}$
4	$\frac{1}{4000}$	$\frac{1}{1600}$	$\frac{1}{1305.9}$	$\frac{1}{1225}$	$\frac{1}{681.8418}$

# Simulation Results

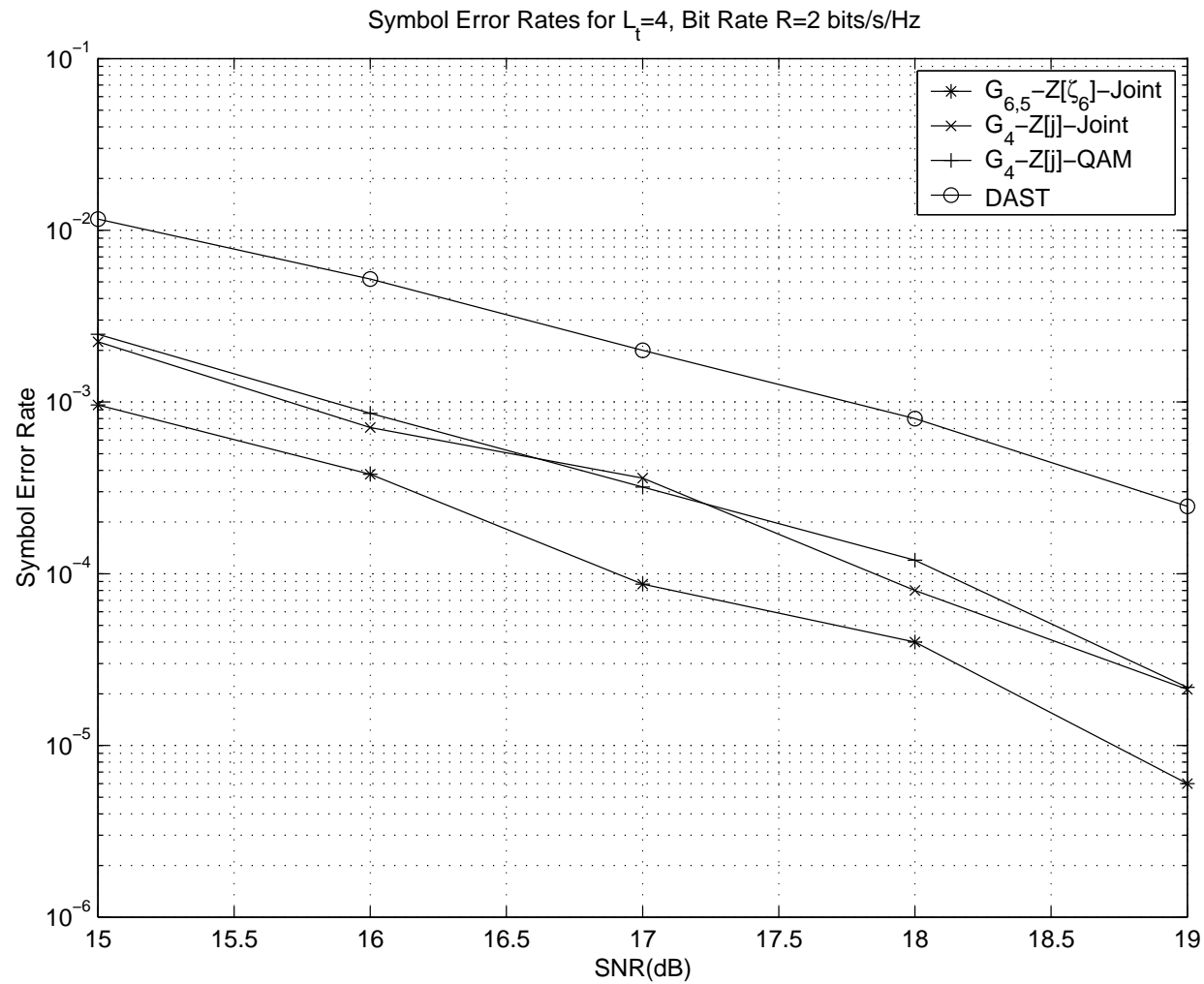


Figure 1: Symbol error rate, information rate  $R = 2$  bits/s/Hz.

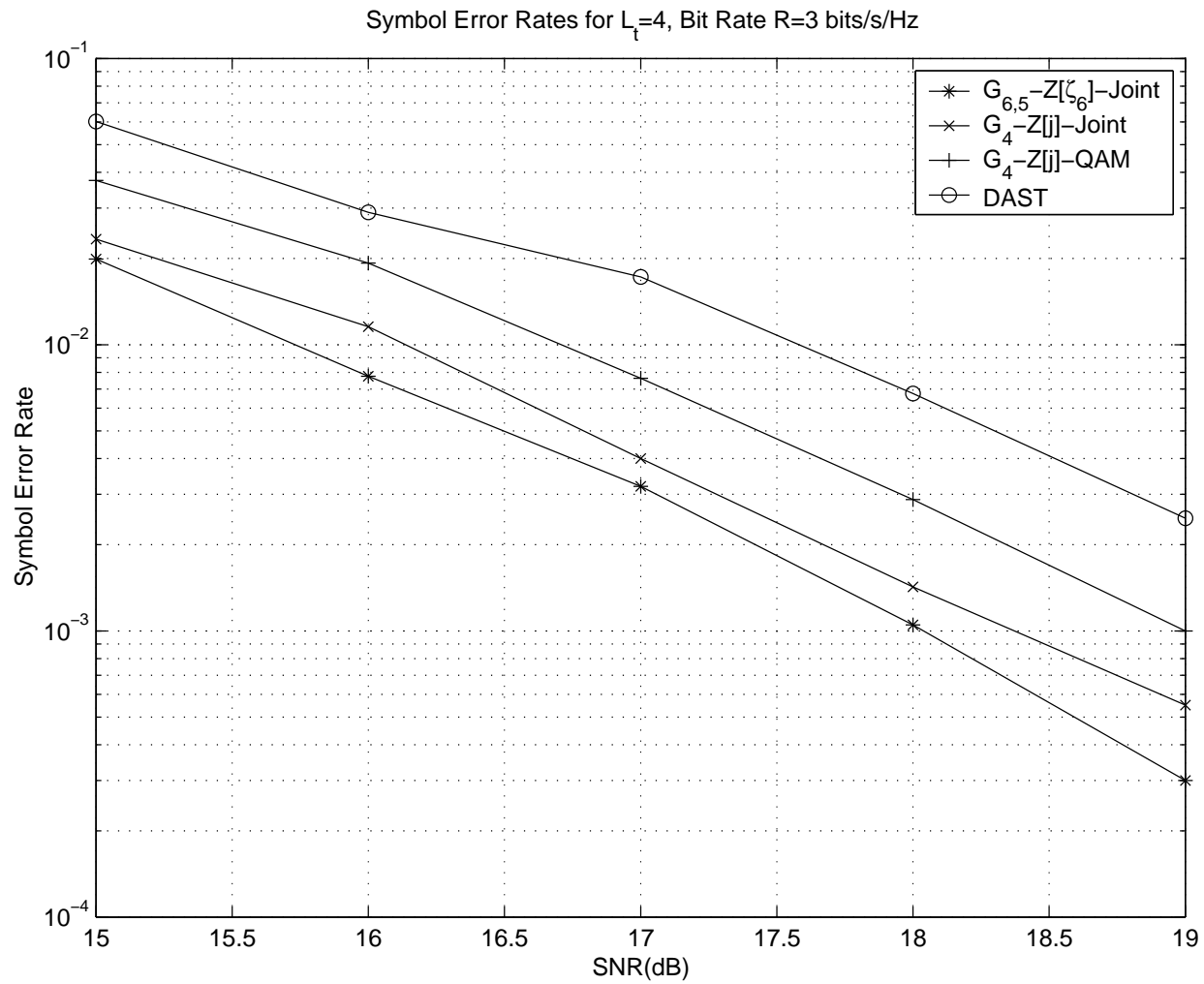


Figure 2: Symbol error rate, information rate  $R = 3$  bits/s/Hz.

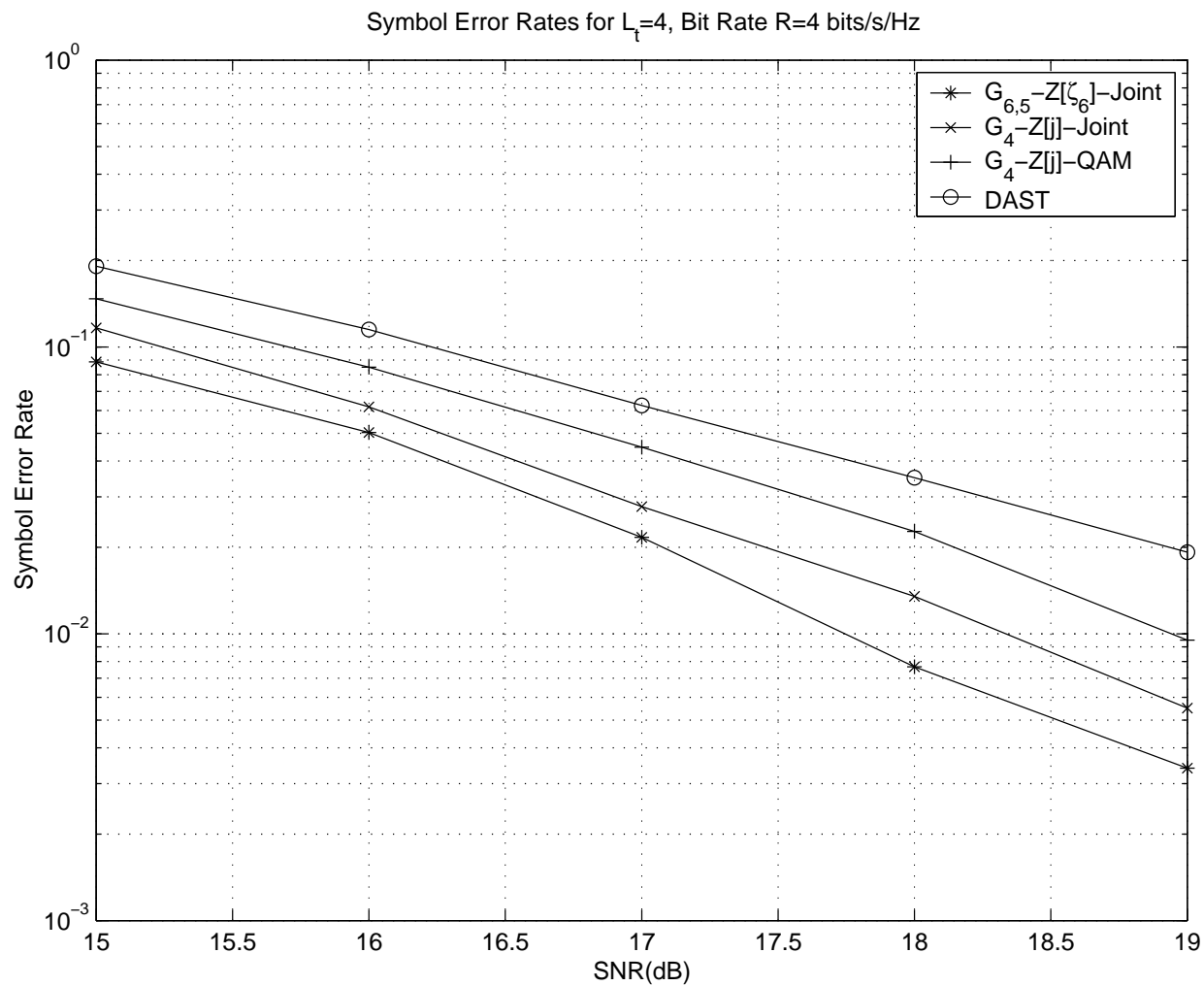


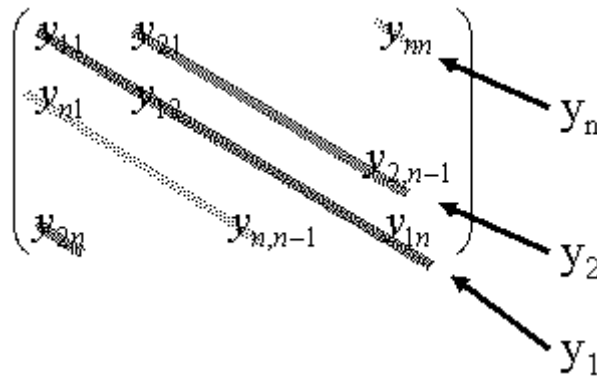
Figure 3: Symbol error rate, information rate  $R = 4$  bits/s/Hz.



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## More Results

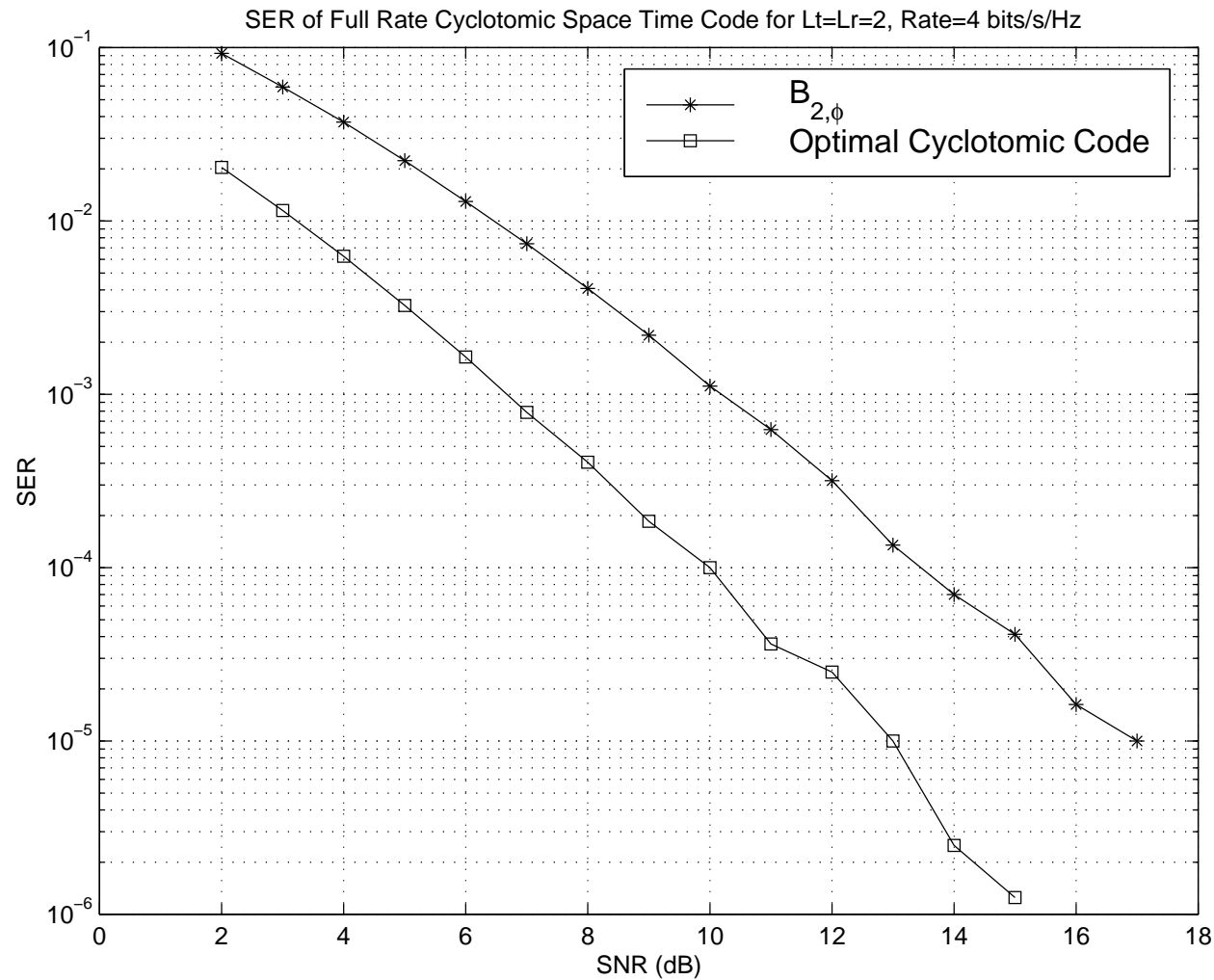
- We have generalized the optimal diagonal cyclotomic space-time codes to multi-layer cyclotomic space-time codes for 2 and 3 transmit antennas



- Obtained some results on super-orthogonal trellis codes from QAM constellations
- Obtained some results on unitary space-time code designs from APSK signals

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- Obtained a new family of recursive space-time trellis codes
  - Obtained closed-form designs of complex orthogonal designs of rates  $(k + 1)/(2k)$  for  $2k$  or  $2k + 1$  transmit antennas.
  - Obtained some new unitary space-time codes with best and best known diversity product diversities by using packing theory.
  - Obtained optimal quasi-orthogonal space-time codes with minimum decoding complexity.
  - Obtained a fast iterative decoding algorithm for lattice based space-time codes based on soft interference cancellation.
  - Applied some of our newly proposed space-time codes into a relay sensor network.

# Simulation Result for Our Optimal Full Rate Full Diversity Code: $R = 4$ bits/s/Hz



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## Future Research

- Space-time code designs beyond cyclotomic rings: for example,  $\mathbb{Z}[\zeta_m, \sqrt{5}]$
- Optimal multi-layer cyclotomic and quadratic space-time codes for more than 3 transmit antennas
- Space-time code designs for relay sensor networks to achieve optimal cooperative diversity
- Recursive space-time trellis code designs with optimal diversity products
- Super quasi orthogonal space-time trellis code designs

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## Conclusion

- Proposed systematic and concrete full diversity cyclotomic lattices and space-time block codes.
- Proposed optimal cyclotomic lattices and space-time block codes by minimizing the mean transmission power when their diversity products are fixed.

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